

## Examples in pari for BSGS, rho, and index calculus

The following is the Pari history with annotations (after the //). I'm skipping lines with typos.

Setting up the group with large prime subgroup:

```
l= 53 // we want a "big" subgroup
```

```
? isprime(2*l+1)
```

```
%7 = 1
```

```
? p=2*l+1
```

```
%8 = 107 // so we'll work in F_107, in the subgroup of order 53
```

```
? Mod(23^2,p) // find a generator by picking a random elt, raising to the power of the cofactor, here just 2
```

```
%9 = Mod(101, 107)
```

```
? g=% // the result was not 1, so g has order 53 and will be our generator
```

```
%10 = Mod(101, 107)
```

```
? znorder(g) // just double checking the order (only for demo puposes, not necessary)
```

```
%11 = 53
```

```
? znorder(Mod(42,p)) //42 is a good target, but is it in the subgroup?
```

```
%12 = 53 // we could also check this by computing 42^53 and getting 1
```

```
? h=Mod(42,p) // yes, it is, so we're taking it as our target
```

```
%14 = Mod(42, 107)
```

```
? m = floor(sqrt(53)) // set up the boundary for BSGS
```

```
%16 = 7
```

```
? g0 =g^0 // begin of baby steps
```

```
%17 = Mod(1, 107)
```

```
? g1 = g^1
```

```
%18 = Mod(101, 107)
```

```
? g2=g^2
```

```
%19 = Mod(36, 107)
```

```
? g3=g2*g
```

```
%20 = Mod(105, 107)
```

```
? g4=g3*g
```

```
%21 = Mod(12, 107)
```

```
? g5=g4*g
```

```
%22 = Mod(35, 107)
```

```
? g6=g5*g // last baby step
```

```

%23 = Mod(4, 107)
? gminv=g^(-7) // compute g^{-7}; normally this would be (g6*g)^{-1}.
%24 = Mod(49, 107)
? h // first giant step, this checks a1 = 0
%25 = Mod(42, 107)
? h*gminv // second giant step, checks a1 = 1
%26 = Mod(25, 107)
? %*gminv // third giant step, checks a1 = 2
%27 = Mod(48, 107)
? % * gminv //fourth giant step, checks a1= 3
%28 = Mod(105, 107) //success, we know this one from the table of BS!

```

```
? a1 = 3
```

```
%29 = 3
```

```
? a0 =3
```

```
%30 = 3
```

```
? a = a0 + m*a1 // putting a together
```

```
%31 = 24
```

```
? g^24 // test
```

```
%32 = Mod(42, 107) // Hurray! This matches the target.
```

### **Pollard rho, schoolbook method.**

```
% g % 3
```

```
%33 = Mod(0, 1) // this does not do what we want: 3 and 107 are coprime, so we get a result mod 1.
```

```
? g
```

```
%34 = Mod(101, 107) // here we see that g carries the Mod 107
```

```
? lift(g) % 3 // lifting is the opposite of reduction, so lift(g) is an integer, so we can reduce that mod 3.
```

```
%35 = 2
```

```
? s = g // start of the slow walk at g
```

```
%36 = Mod(101, 107)
```

```
? f = g // start of the fast walk at g
```

```
%37 = Mod(101, 107)
```

```
? s= s^2 // 101 is 2 mod 3, so we square
```

```
%38 = Mod(36, 107)
```

```
? bs = 1 // we need to set the counters
```

```
%39 = 1
```

```
? bs = 2*bs // we need to update the counters
```

```
%40 = 2
? cs = 0 // setting counters for power of h
%41 = 0
? cs = 2*cs //updating
%42 = 2
? bf = 1 //same for fast walk
%43 = 1
? cf = 0
%44 = 0
? f = f^2 // fast walk starts; we know this step already, but we need to get all variables right
%47 = Mod(36, 107)
? bf = 2*bf //updates
%48 = 2
? cf = 2*cf
%49 = 0
? lift(f) % 3 // second part of the fast step is *g
%50 = 0
? f = f*g
%51 = Mod(105, 107)
? bf = bf +1
%52 = 3
? cf = cf
%53 = 0 // done with 1 fast step; will skip updates to c until they are non-zero
? lift(s) % 3 // begin of the 2nd slow step
%54 = 0
? s = s*g
%55 = Mod(105, 107)
? bs = bs +1
%56 = 3
? lift(f) %3 // begin of the 2nd fast step
%57 = 0
? f = f*g
%58 = Mod(12, 107)
? bf = bf +1
```

%59 = 4

? lift(f) %3 // and the second half of it

%60 = 0

? f = f\*g

%61 = Mod(35, 107)

? bf = bf +1

%66 = 5

? lift(s) % 3 // 3<sup>rd</sup> slow step

%67 = 0

? s = g\*s

%68 = Mod(12, 107)

? bs = bs +1

%69 = 4

? lift(f) % 3 // 3<sup>rd</sup> fast step

%70 = 2

? f = f^2

%71 = Mod(48, 107)

? bf = 2\*bf

%72 = 10

? f=f\*g //I can read off that 48 %3 = 0

%73 = Mod(33, 107)

? bf = bf+1

%74 = 11

? s= s\*g //begin of 4<sup>th</sup> step

%75 = Mod(35, 107)

? /// here my laptop was going crazy and I didn't see the result, so I retyped it;

? s = s\*g // this is not the next step, as 35 %3 = 2 and not 0; so I left the path, oops!

%76 = Mod(4, 107) // this is not there, I should still be at 35

? bs = bs+1

%77 = 5

? f = f\*g // 4<sup>th</sup> fast step

%78 = Mod(16, 107)

? bf = bf+1

%79 = 12

? f = f\*h // finally a step involving h!

%80 = Mod(30, 107)

? cf =cf +1

%81 = 1 // this is when I noticed that I had done an extra step, but

*Try to finish this yourself; you should find a=24 again*

### **Index calculus attacker**

Trying to solve the same DLP, this time with index calculus.

107 is small, so  $F=\{2,3,5\}$ .

The running time depends on the smoothness probability of numbers in  $[0,p-1]$ .

? g^12 // we pick some random exponent of g

%89 = Mod(16, 107)

? factor(lift(%)) // then factor the integer representation of the result

%90 =

[2 4] // OK, this is quite unusual! I expect a mix of primes, but hey, we get the DL of 2 with base g

? g // reminder of what g is

%91 = Mod(101, 107)

// this equation means that  $12 = 4 * \log_g(2)$ , put  $\lg_2 = \log_g(2)$

? lg2=3

%92 = 3

? g^17 // another random exponent

%93 = Mod(25, 107) // again, too nice

? lg5=Mod(17/2,53) // note that we're working with exponents, so those are mod 17

? g^23

%95 = Mod(100, 107) // again too nice and we know lg2 and lg5 already

? g^27

%96 = Mod(23, 107) // more normal bad luck

? g^42

%97 = Mod(13, 107) // more normal bad luck

? g^18

%98 = Mod(64, 107) // too nice, but nothing new

? g^21

%99 = Mod(86, 107)

? factor(lift(%)) // this is how you would implement this

```
%100 =  
[ 2 1]  
[43 1]  
? g^29  
%101 = Mod(79, 107) // nope  
  
? factor(lift(%))  
%102 =  
[79 1]  
? g^37  
%103 = Mod(34, 107)  
? factor(lift(%))  
%104 =  
[ 2 1]  
[17 1]  
? g^random(53) // somebody was concerned about me depleting my brain of random numbers  
%105 = Mod(49, 107)  
? factor(lift(%))  
%106 =  
[7 2] // oh, nice, but 7 is not in our factor base, so ignore this  
? g^random(53)  
%107 = Mod(14, 107) // same  
[[snip]]  
? g^random(53)  
%114 = Mod(81, 107)  
? factor(lift(%)) // oh, nice! We would know all DLs for the factor basis  
// alas I don't know what random exponent that was, oops!  
%115 =  
[3 4]  
? b= random(53) // let's do it the safe way  
%116 = 44  
? g^b  
%118 = Mod(40, 107)  
? factor(lift(%))
```

```

%119 = // nothing new here
[2 3]
[5 1]
? b= random(53) // let's just move to the second phase; maybe 3 isn't even in the right subgroup
%121 = 36
? h*g^b
%122 = Mod(83, 107)
? factor(lift(%)) //nope
%123 =
[83 1]
? factor(lift(h)) // also should try h itself, but no
%124 =
[2 1]
[3 1]
[7 1]
? b= random(53)
%125 = 23 //great randomness!
? h*g^b
%126 = Mod(27, 107) // argh, now we would need dg3
? b= random(53)
%128 = 7
? h*g^b
%129 = Mod(62, 107) // nope
? b= random(53)
%130 = 3
? h*g^b
%131 = Mod(23, 107) //nope
? b= random(53)
%132 = 14
? h*g^b
%133 = Mod(10, 107) // yeah! This factors of the factor base and we know both pieces
? factor(lift(%))
%134 =
[2 1]

```

[5 1]

$A = \text{Mod}(-14 + 1 \cdot \lg 2 + 1 \cdot \lg 5, 53)$

$\%136 = \text{Mod}(24, 53)$

$g^A * g^{14} - g^{(\lg 2 * 1)} * g^{(\text{lift}(\lg 5) * 1)}$  // this is how we found it,  $\text{lift}(\lg 5)$  because it was mod 53

$\%139 = \text{Mod}(0, 107)$