

# Coppersmith / Howgrave-Graham and LLL

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# Theorem by Howgrave-Graham

Let  $g(x) = \sum_{i=0}^{d-1} g_i x^i \in \mathbf{Z}[x]$  of  $\deg(g) = d - 1$ .

Let  $m, b \in \mathbf{Z}_{>0}$

If

1.  $g(x_0) \equiv 0 \pmod{b^m}$  with  $|x_0| \leq X$ ,
2.  $\|g(xX)\| \leq b^m / \sqrt{d}$

then  $g(x_0) = 0$  over  $\mathbf{Z}$ .

Here  $\|g(xX)\| = \sqrt{g_0^2 + g_1^2 X^2 + \dots + g_{d-1}^2 X^{2(d-1)}}$  is the Euclidean norm of the coefficient vector of  $g(xX)$ .

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If  $g(x_0) \in b^m \mathbf{Z}$ ,  $|x_0| < X$ ,  $|g_0| \leq b^m/d$ ,  $|g_1 X| \leq b^m/d$ ,  
 $|g_2 X^2| \leq b^m/d, \dots, |g_{d-1} X^{d-1}| \leq b^m/d$  then  $g(x_0) = 0$ .

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Now just square, sum up and take the square root on both sides:

$$\sum_{i=0}^{d-1} g_i^2 X^{2i} \leq \sum_{i=0}^{d-1} b^{2m}/d^2 = b^{2m}/d.$$



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If that's too restrictive we can expand to

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## What to look for and how to find it?

All of these attacks start by finding some polynomial  $\deg(f) = t$  for which we have a root modulo  $b^m$  is interesting. Let  $\deg(f) = t$  and let  $|x_0| \leq X$  for some known  $X$ .

To find

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we will use LLL, which builds integer linear combinations of the input rows of a matrix. It returns a vector that is short in the Euclidean norm. (Hence we wanted that in the Howgrave-Graham theorem).

We set up a system of equations in the coefficient vectors, one row per option.  $b^m \mathbf{Z}$  turns into coefficient  $b^m$  at the  $x^0$  column etc.

For Howgrave-Graham we need to scale the column of  $x^s$  by  $X^s$ . So we get

$$\begin{pmatrix} X & a \\ 0 & n \end{pmatrix}$$

for knowing part of  $p$ .



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$$\begin{pmatrix} X^2 & aX & 0 \\ 0 & X & a \\ 0 & 0 & n \end{pmatrix} \stackrel{?}{=} \begin{pmatrix} X^2 & 2aX & a^2 \\ 0 & X & a \\ 0 & 0 & n \end{pmatrix}$$

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# LLL

Due to Lenstra, Lensra, and Lovász, 1982.

- ▶ On input a set of vectors  $\{v_1, v_2, \dots, v_d\}$  output a short vector  $v'_1$  so that  $v'_1 = \sum a_i v_i$  for some  $a_i \in \mathbf{Z}$ .

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- ▶ LLL outputs  $d$  vectors which are shorter and more orthogonal. Each vector is an integer linear combination of the inputs.
- ▶ LLL uses many elements from Gram-Schmidt orthogonalization:
  - ▶ for  $j = 1$  to  $d$
  - ▶ for  $i = 1$  to  $j - 1$
  - ▶ 
$$\mu_{ij} = \frac{\langle v_i^*, v_j \rangle}{\langle v_i^*, v_i^* \rangle}$$
  - ▶ 
$$v_j^* = v_j - \sum_{i=1}^{j-1} \mu_{ij} v_i^*$$
- ▶ Note that the  $\mu_{ij}$  are not integers, so the  $v_j^*$  are not in the lattice.
- ▶ A lattice basis is LLL reduced for parameter  $0.25 < \delta < 1$  if
  - ▶  $|\mu_{ij}| \leq 0.5$  for all  $1 \leq j < i \leq d$ ,
  - ▶  $(\delta - \mu_{i-1,i}^2) \|v_{i-1}^*\|^2 \leq \|v_i^*\|^2$ .
- ▶ This guarantees  $\|v_1\| \leq (2/\sqrt{4\delta - 1})^{(d-1)/2} \det(L)^{1/d}$ , where  $\det(L)$  is the determinant of the lattice.

## LLL algorithm (from Cohen, GTM 138, transposed)

Input: Basis  $\{v_1, v_2, \dots, v_d\}$  of lattice  $L$ ,  $0.25 < \delta < 1$

Output: LLL reduced basis for  $L$  with parameter  $\delta$

1.  $k \leftarrow 2$ ,  $k_{\max} \leftarrow 1$ ,  $v_1^* \leftarrow v_1$ ,  $V_1 = \langle v_1, v_1 \rangle$
2. if  $k \leq k_{\max}$  go to step 3  
else  $k_{\max} \leftarrow k$ ,  $v_k^* \leftarrow v_k$ , for  $j = 1$  to  $k - 1$ 
  - ▶ put  $\mu_{jk} \leftarrow \langle v_j^*, v_k \rangle / V_j$  and  $v_k^* \leftarrow v_k^* - \mu_{jk} v_j^*$ $V_k = \langle v_k, v_k \rangle$
3. Execute RED( $k, k - 1$ ). If  $(\delta - \mu_{i-1,i}^2) V_{k-1} > V_k$  execute SWAP( $k$ ) and  $k \leftarrow \max\{2, k - 1\}$ ; else for  $j = k - 2$  down to 1 execute RED( $k, j$ ) and  $k \leftarrow k + 1$ .
4. If  $k \leq d$  go to step 2; else output basis  $\{v_1, v_2, \dots, v_d\}$ .
  - ▶ RED( $k, j$ ): If  $|\mu_{jk}| \leq 0.5$  return; else  $q \leftarrow \lfloor \mu_{jk} \rfloor$ ,  $v_k \leftarrow v_k - qv_j$ ,  $\mu_{jk} \leftarrow \mu_{jk} - q$ , for  $i = 1$  to  $j - 1$  put  $\mu_{ik} \leftarrow \mu_{ik} - q\mu_{ij}$  and return.
  - ▶ SWAP( $k$ ): Swap  $v_k$  and  $v_{k-1}$ . If  $k > 2$  for  $j = 1$  to  $k - 2$  swap  $\mu_{jk}$  and  $\mu_{j,k-1}$  and update all variables to match (see p.88 in Cohen)

For a nice visualization see pages 61–66 of

<http://thijs.com/docs/lec1.pdf>.

(Animations only work in acroread.)