

Thanks a lot to the people who took notes!

Examples how RSA works in practice

package name 'pari-gp'

Pari-gp is a general computer algebra system; this is what Tanja will provide on the laptops if the exam takes place in person. It has all the number theory functions we need and very large precision.

https://pari.math.u-bordeaux.fr/dochtm/html/Arithmetic_functions.html#isprime

Check if p and q are prime (Fermat's primality test):

$$a^{p-1} \equiv 1 \pmod{p} \text{ for all } a \text{ with } \gcd(a,p) = 1$$

Example in practice:

$$p = 103$$

$$q = 107$$

Check base $a = 2, 3, 5$

Primality test is true \rightarrow p probably prime or a Carmichael number

Primality test is false for some $a \rightarrow$ p is certainly not prime

Here: all three come back as " p probably prime or a Carmichael number"

We can also check that 103 is prime by checking for prime divisors up to \sqrt{p} , in this case there is none, so 103 is prime.

in gp type

`isprime(103)`

which answers 1 to say that 1 is prime (check out the documentation for different options)

Next compute n :

$$n = p \cdot q = 11021$$

Next compute ϕ :

$$\phi = (p-1)(q-1) = 10812$$

Now, for exponentiation, can we use $e = 3$? Basically no. (Why? because ϕ was divisible by 3)

Use $e = 5$

Compute bezout (=extended Euclidean algorithm):

`bezout(5, 10812)`

$$= (4325, -2, 1)$$

meaning that

$$4325 \cdot 5 - 2 \cdot 10812 = 1$$

`bezout` is Pari-GP's command for XGCD.

To compute XGCD by hand:

Check quotient and remainder

$$10812/5.0 = 2162.4000000$$

$$q = 2162$$

$$r = 2$$

then

$5/2.0 = 2.5$, so the next quotient is 2

$$r = 1$$

$$d = 2162 \cdot 2 + 1$$

Using the algorithm which will appear on the exercise sheet (use that if you don't have your favorite way to compute XGCD -- or to teach it to your calculator):

$$\begin{array}{r} 10812 \quad 1 \quad 0 \\ \quad 5 \quad 0 \quad 1 \\ \quad 2 \quad 1 \quad -2162 \\ \quad 1 \quad -2 \quad 4325 \end{array}$$

Check whether $e \cdot d \bmod(\phi) = 1$

Yup $\text{Mod}(5 \cdot 4325, \phi) = \text{Mod}(1, \phi)$

Difference mod operator and %

"Mod" operator gives remainder and modulus, residue class

"%" only gives remainder, the integer

Key generation

Create public key: (n,e)

Private key: (n,d)

pgpdump to display pgp key

Format specified in RFC 4880

$e = 01\ 00\ 01$ (binary) means $2^{16} + 1 = 65537$

Showing Secret/private key in pgpdump:

displays n, e, d, p, q and u in hexadecimal

Recall $u = p^{-1} \pmod{q}$

Alice's public key (n,e) = (11021,5)

Alice's private key (n,d) = (11021,4325)

Somebody wants to encrypt a message to Alice

$$m = 1234 < n$$

$$m^5 = \dots$$

Use square and multiply:

$$5 = 2^2 + 1 = [101]_{\{2\}}$$

$n=11021$
 $\text{Mod}(m^2, 11021) = \text{Mod}(1858, 11021)$
 Bit is 0, so no mult
 $\text{Mod}(1858, 11021)^2 = \text{Mod}(2591, 11021)$
 Bottom bit = 1, so multiply by m
 $\text{Mod}(2591, 11021) * \text{Mod}(m, n) = \text{Mod}(1204, 11021)$

(Use %N to refer to output N in pari-gp)

(Pari reduces along the way ... and so should you!)

CRT method:

Example:

Bob has $n = 164063$, $e = 17$ and receives

$c = 6215$

He used $p = 359$ and $q = 457$ and got

$d = 57617$

$dp = d \% (p-1)$ (Exponent gets smaller) = 337

$dq = d \% (q-1) = 161$

$cp = c \% p = 112$

$cq = c \% q = 274$

$mp = \text{Mod}(cp^{dp}, p) = \text{Mod}(89, 359)$

$mq = \text{Mod}(cq^{dq}, q) = \text{Mod}(172, 457)$ (mp & mq are useful when calling chinese(), i.e 2 inputs instead of 4)

$\text{chinese}(mp, mq) = \text{Mod}(75120, 164063)$

chinese=CRT in pari-gp), the inputs to chinese are of the form $\text{Mod}(a,b)$; which matches the above

Is the CRT calculation a bottleneck? No, $u = p^{-1} \text{ mod } q$ is included in key. Then

$m = mp + u * p * (mp - mq) \text{ mod } n$

(you can compute this by first computing $p * (mp - mq) \% q$ and then multiplying the result of this by u and adding mp , this typically means that no reduction mod n is needed.

Remarks:

- 1) important to reduce $c \text{ mod } p$ and $\text{mod } q$, and to reduce $d \text{ mod } p-1$ and $\text{mod } q-1$
- 2) exponents are $\text{mod } \phi(p)$ etc., needs that base and p are coprime

If an attacker disturbs Bob in the RSA-CRT computation so that he accidentally computes

$mp = \text{Mod}(165, 359)$ and then gets plaintext

$m' = \text{chinese}(mp, mq) = \text{Mod}(129046, 164063)$

as attacker I know that the correct plaintext is 75120, what can you learn?

Hint:

$\text{Mod}(75120 - 129046, q) = \text{Mod}(0, 457)$

$\text{Mod}(75120 - 129046, p) = \text{Mod}(283, 359)$

Considering the homework sheet:

The homework sheet will include algorithms for XGCD and CRT if you are not sure how to do these (or

how to teach them to you computer)

Submit homework by email

Postscript

I had messed up in copying the numbers. what I wanted to give you with CRT is

$c = 66215$, $d = 57617$, $p = 359$, $q = 457$.

Figure out what the plaintext was.