

Extended Euclidean algorithm (XGCD)

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2WF80: Introduction to Cryptology

Euclidean algorithm and gcd

- ▶ The Euclidean algorithm computes the gcd of two numbers

$$d = \text{gcd}(m, n)$$

in time polynomial in $\log_2(\max\{m, n\})$.

- ▶ This is much faster than factoring m and n .
- ▶ Each step computes the quotient and remainder of two integers, starting with $m = q_1 \cdot n + r_1$, followed by $n = q_2 r_1 + r_2$,
 $r_1 = q_3 r_2 + r_3, r_2 = q_4 r_3 + r_4, \dots$
The algorithm stops when $r_i = 0$ and outputs $d = r_{i-1}$ as the gcd.

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$$d = \gcd(m, n) = am + bn,$$

and $|a| < n, |b| < m$.

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- ▶ Can compute a, b by reversing steps above, starting with

$$r_{i-1} = r_{i-3} - q_{i-1}r_{i-2} = r_{i-3} - q_{i-1}(r_{i-4} - q_{i-2}r_{i-3}) = \dots = am + bn$$

Extended Euclidean algorithm

Input $m, n \in \mathbb{N}$

Output $d \in \mathbb{N}, a, b \in \mathbb{Z}$ with $d = am + bn$

1. $v \leftarrow [m, 1, 0]$
2. $w \leftarrow [n, 0, 1]$
3. while $w[0] \neq 0$
 - 3.1 $x \leftarrow v - (v[0] \text{ div } w[0]) w$
 - 3.2 $v \leftarrow w$
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4. $d \leftarrow v[0], a \leftarrow v[1], b \leftarrow v[2]$
5. return d, a, b

Extended Euclidean algorithm

Input 312, 213

$$\begin{bmatrix} 312, & 1, & 0 \\ 213, & 0, & 1 \end{bmatrix} \quad q = 1$$

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$d = 3, a = 28, b = -41$

indeed

$$28 \cdot 312 - 41 \cdot 213 = 3.$$

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At every step, $v[0] = v[1]m + v[2]n$.

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XGCD for inversion

- ▶ On input m, n , XGCD computes d, a, b with

$$d = am + bn.$$

- ▶ An integer m is invertible modulo n if it is co-prime to n , i.e., if $\gcd(m, n) = 1$.
- ▶ XGCD is an efficient way to compute modular inverses:

$$1 = am + bn \quad \Rightarrow$$

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In the example

$$28 \cdot 312 - 41 \cdot 213 = 3.$$

Thus 312 and 213 are not co prime.