Privacy Preserving Protocols Workshop on Cryptography for the Internet of Things

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Introduction

└─ Cryptography in Daily Life

RFID







└─ Cryptography in Daily Life

Car Keys







Cryptography in Daily Life

Access Control



Product Tracking





2 Privacy Models

Protocol Analysis Provable Security (Privacy) Privacy Model Insider Attacks Requirements

- 3 Lightweight Cryptography
- **4** Existing Protocols
- 5 Protocol Design Design
- 6 Conclusions and Future Perspectives

Why?



Industrial espionage

Why?



User privacy

Why?



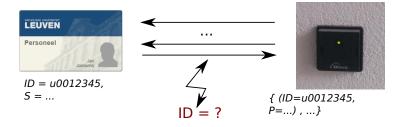
User privacy

Why?

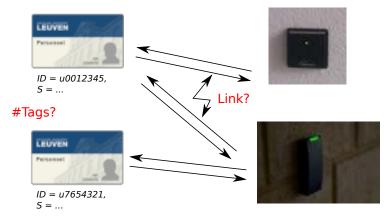


Wireless Gun

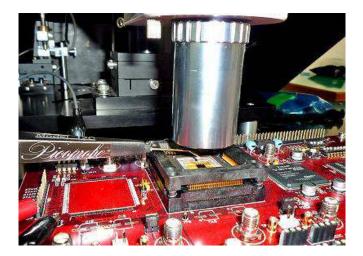
RFID Privacy: goals



RFID Privacy: goals



Corrupting Tags



Different Privacy Solutions

- Protocol Level Privacy
- Kill Command
- Destroy Tag
- Shielding
- (Read Range Reduction)
- **.**...

Threat Analysis / Requirements

		Privacy	
		Low	High
Security	Low	Supply Chain	Public Transport
S	High	Car Keys	Payments Access Control Passports



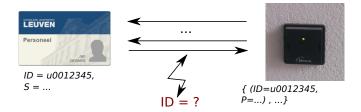
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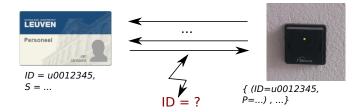
Protocol Analysis



Properties:

- Security
- Privacy: untraceability
- Allow corruption

Protocol Analysis



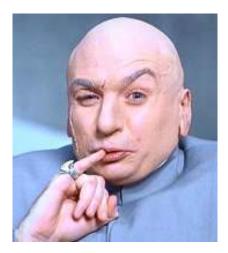
Results

Many published protocols broken:

 \Rightarrow Lack of formal proofs!

Provable Security (Privacy)

Provable Security (Privacy)



Provable Security (Privacy)

Provable Security (Privacy)

Adversary







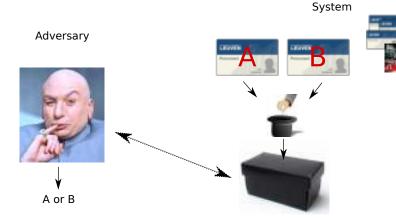
System



Adversary wins if ...

Privacy Preserving Protocols
Privacy Models
Privacy Model

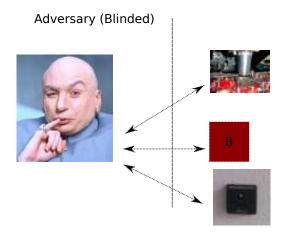
Juels-Weis model (2005)



Adversary wins if output is correct tag.

Vaudenay model (2007)

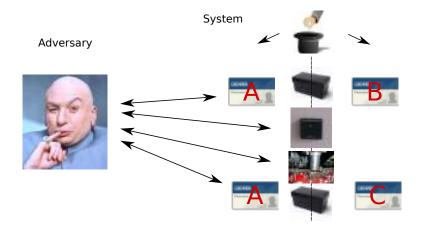
System

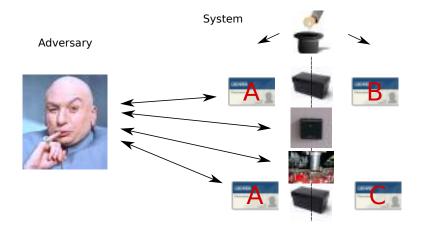


Adversary wins if output is true and not trivial

Design goals:

- Strong adversary: can always corrupt
- Solve issues with wide strong privacy
- Model 'reality'
- Easy to use





Adversary wins if random bit is guessed correctly.

New Features:

- corruption \rightarrow on *real* tag
- wide strong privacy



Features (reused):

- Virtual tag handles
- Indistinguishability based
- Single random bit for entire system

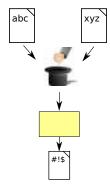


Indistinguishability

Encryption:

- RO
- IND-CPA
- IND-CCA
- IND-CCA2

...



Privacy-models:

- Juels-Weis
- Vaudenay
- Hermans *et al.*

Indistinguishability

Encryption:

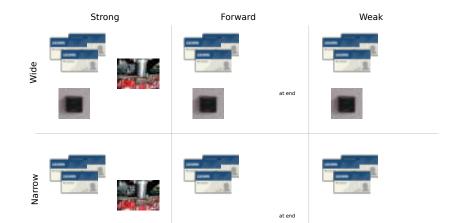
- RO
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Privacy-models:

- Juels-Weis
- Vaudenay
- Hermans et al.

Privacy Levels



Privacy Requirements

Privacy Level	Application
Narrow Weak	Supply Chain
Narrow Forward	Smart Products
Wide Weak	Car Keys
Wide Forward	Payments Access Tokens Passports Public Transport

Insider Attacks

System

Adversary



Insider Tag





— Requirements

Privacy Requirements

Privacy Level	Application
Narrow Weak	Supply Chain
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Wide Forward + Insider	Payments Access Tokens Passports Public Transport

— Requirements

Privacy Requirements

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1 RFID Privacy Requirement

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Lightweight Devices















Lightweight Cryptography?



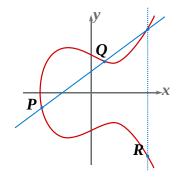
Limits:

- Area (€€€)
- Time
- Power
- Energy

Typical Ingredients for Protocols

Primitive	Status	
RNG	OK?	
Key Update	???	
Block Cipher	OK	
Hash Function	OK	
ECC	OK	
\sum	???	

Lightweight Elliptic Curve Cryptography



Implementation [LBSV10]:

- Area (14.5 kGE)
- Time (85 ms)
- Power (13.8 µW)
- Energy (1.18 µ J)

RFID Privacy Requirement

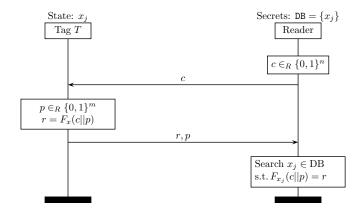
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PRF (Block cipher) based [ISO/IEC 9798-2]



Privacy Wide-Weak

Symmetric Key and Efficiency

Damgård-Pedersen '08:

- Independent keys: inefficient O(n)
- Correlated keys:
 - efficient $O(\log(n))$
 - privacy loss

Symmetric Key and Efficiency

Damgård-Pedersen '08:

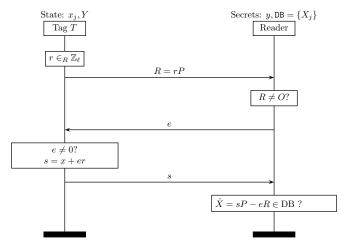
- Independent keys: inefficient O(n)
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Key Updating

- Higher Privacy Level (narrow forward)
- Desynchronization Attacks / Efficiency Problems
- Implementation cost?

Privacy Preserving Protocols

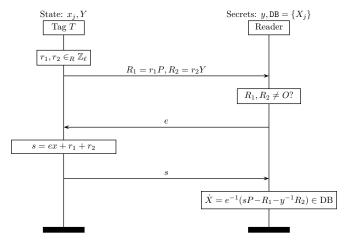
EC Schnorr Protocol



Privacy

None

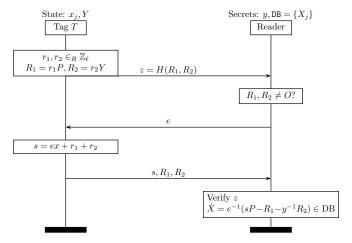
Randomized Schnorr [BCI08]



Privacy

Narrow Strong

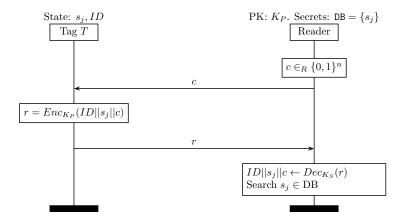
Randomized Hash GPS [BCI09]



Privacy

Narrow Strong and Wide Forward

IND-CCA2 Encryption [Vau07]



Privacy Wide Strong

Protocol	Privacy	Ins.	Ext. Snd.	Operations
Schnorr	no	no	yes	1 EC mult
Randomized Schnorr	narrow-strong	no	yes	2 EC mult
Rand. Hashed GPS	narrow-strong wide-forward	no	yes	2 EC mult 1 hash

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Hash ElGamal	wide-strong	yes	no	2 EC mult 1 hash 1 MAC



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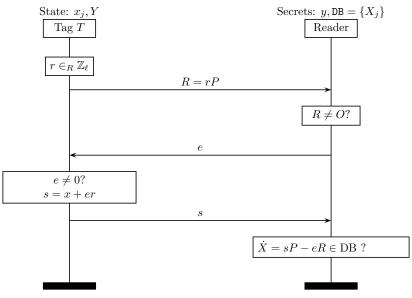
New Protocol [Peeters, Hermans 2012]

Design protocol:

- Correct
- Extended soundness
- (At least) Wide Forward + Insider privacy
- Efficient

Privacy Preserving Protocols

EC Schnorr Protocol



Oracle Diffie-Hellman Assumption

$$(A = aP, B = bP, abP) \sim (A = aP, B = bP, rP)$$

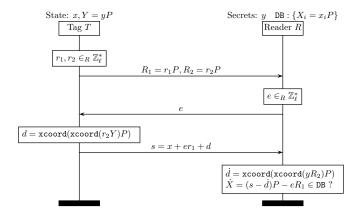
with extra $\mathcal{O}(Z) := \operatorname{xcoord}(bZ)P$.

X Logarithm

 $xcoord(rP)P \sim r'P$

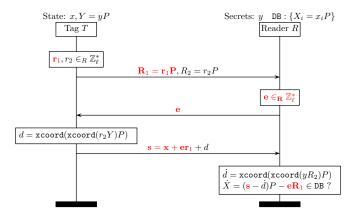
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Privacy Preserving Protocols
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New Protocol



Privacy Preserving Protocols

New Protocol - Extended Soundness

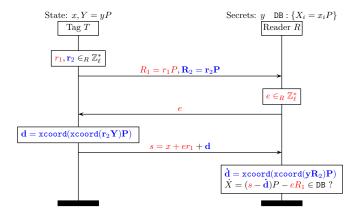


Extended Soundness

Schnorr protocol \Rightarrow extended soundness (OMDL assumption)

Privacy Preserving Protocols

New Protocol - Privacy



Performance

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Performance

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Our Protocol - optimised version	wide-forward-insider wide-forward-insider	yes yes	yes yes	4 EC mult 2 EC mult

Summary

- Overview RFID Privacy Models & Privacy Levels
- Implementation Aspects
- RFID Protocols
- New Private & Efficient RFID Protocol

Future Perspectives

Privacy models

- 'Fair' comparison
- Restrictions on tag corruption
- Simulatability vs indistinguishability

Protocols

- New applications
- Other primitives \rightarrow feasible?
- Analyze underlying assumptions (DDH-variants)



Privacy Preserving Protocols Conclusions and Future Perspectives

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